

# CS530

## Key Management & Distribution Issues (Part 2)

Bill Cheng

<http://merlot.usc.edu/cs530-s10>

Copyright © William C. Cheng



2

## Key Management & Distribution Issues

### Key Management & Distribution Issues

#### Practical issues

↳ how to carry them

○ passwords vs. disks vs. smartcards

↳ where do they stay, where do they go

↳ how many do you have

↳ how do you get them to begin with

#### Classes of crypto

↳ which type is right for your application

#### Who needs strong secrets?

#### How do you recover from exposed keys?

#### Miscellaneous issues

↳ security architectures

3



Copyright © William C. Cheng

## Key Management Overview

### Key management is where much security weakness lies ↳ what types of keys to use for a system and how to choose keys

○ want large **key entropy** (amount of randomness in keys)

○ nobody uses rot13 – (inverse is itself)

○ example of weak protocol: WEP

○ really short keys: PIN

○ verifiable plaintext attacks

Ex: Does this look like English?

If plaintext contains a checksum, great! Let's automate the attack!

○ known plaintext attacks

Ex: precomputed dictionary attack

↳ need to **salt** the password (then can only use

dictionary attack)

Copyright © William C. Cheng

4



## Key Management Overview (Cont...)

- ↳ where do you store the keys?
  - floppy disks, USB harddrives (can be encrypted)
  - smartcard
    - key never leaves card
    - not vulnerable to even keyboard sniffer
    - not popular in US, probably because high costs (cost of cards + cost of infrastructure)
    - variety of smartcards: tamper proof, tamper resistant, tamper evident (tamper evident is good enough for end users)
    - post-it note?

5



## Key Management Overview (Cont...)

- ↳ Classes of crypto
  - one-time pad (truly random)
    - most secure, not vulnerable to attacks
    - if pseudo-random number generator used, must have large IV
    - problem: key size must be as large as data size
    - limited applications
    - Ex: submarines
      - visual cryptography (next page)

6



## Key Management Overview (Cont...)

### how do you communicate about keys (key distribution)?

○ conventional: KDC

○ single key shared by both parties

○ generate and distribute keys

○ bind names to shared keys

○ public key: CA

○ public key published to the world

○ private key known only by owner

○ sign bindings of keys to names (protects integrity)

○ verifiable by multiple parties

○ third party certifies or distributes keys

○ certification infrastructure

○ authentication

Copyright © William C. Cheng

7



## Visual Cryptography

Invented by Naor and Shamir (presented at EUROCRYPT '94)  
see [Doug Stinson's Visual Cryptography Page](#)



input pixels	key	output shares 1	output shares 2	merged shares
□	p=0.5	□	□	□
□	p=0.5	□	■	□
■	p=0.5	■	□	■
■	p=0.5	■	■	■

2x data expansion

- application: secret splitting
- **perfect secrecy** (just like a one-time pad)

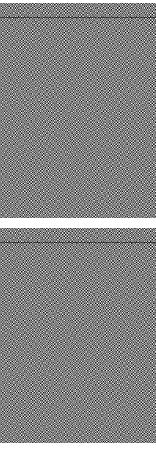


Copyright © William C. Cheng

## Visual Cryptography Example



Ex:



Copyright © William C. Cheng

## Visual Cryptography Example (Cont...)



Ex:



- original image

SEND  
MORE  
MONEY



Copyright © William C. Cheng

CSCI 530, Spring 2010

CSCI 530, Spring 2010

## Grey or Color Image



- If a pixel value is not just black/white
  - grey images - real value between 0 (black) and 1 (white)
  - color images - RGB, real value between 0 and 1 in each component color
  - in both cases, can approximate with pure black and white values



- Two basic approaches
  - **thresholding** - e.g., replace value by 1 if intensity  $\geq 0.5$  and replace value by 0 if intensity  $< 0.5$
  - **error diffusion** -- start with thresholding, carry error into the next pixel



Copyright © William C. Cheng

CSCI 530, Spring 2010

CSCI 530, Spring 2010

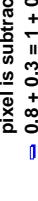
## Key Management Overview (Cont...)



- Who needs strong secrets anyway? (sometimes, secrets are not needed, what is really needed is **integrity of association**)
  - users?
    - need to prove identity
      - start with something not that confidential (SSN, mother's maiden name)
    - servers?
      - private key is usually sitting on the server!
        - not well protected
        - should probably put it on a smartcard



## Error Diffusion



- Error diffusion
  - if a pixel value is 0.8, approximate it with 1, the difference (0.2) is the error
  - if the next pixel value is 0.3, the error in the previous pixel is subtracted, the resulting pixel value is 0.1
    - $0.8 + 0.3 = 1 + 0.1$
  - 0.1 is approximated by 0, the new error is 0.1
- keep going



original



Copyright © William C. Cheng



thresholding



Copyright © William C. Cheng



error diffusion

CSCI 530, Spring 2010

CSCI 530, Spring 2010



Copyright © William C. Cheng



## Key Management Overview (Cont...)

- Who needs strong secrets anyway? (sometimes, secrets are not needed, what is really needed is **integrity of association**)
  - users?
    - need to prove identity
      - start with something not that confidential (SSN, mother's maiden name)
    - servers?
      - private key is usually sitting on the server!
        - not well protected
        - should probably put it on a smartcard



12

CSCI 530, Spring 2010

CSCI 530, Spring 2010



## Key Management Overview (Cont...)

- Who needs strong secrets anyway? (sometimes, secrets are not needed, what is really needed is **integrity of association**)
  - users?
    - need to prove identity
      - start with something not that confidential (SSN, mother's maiden name)
    - servers?
      - private key is usually sitting on the server!
        - not well protected
        - should probably put it on a smartcard



12

CSCI 530, Spring 2010

CSCI 530, Spring 2010

## Key Management Overview (Cont...)

- ↳ Who needs strong secrets anyway? (cont...)
  - DRM? (Digital Rights Management)
    - does it really work? (e.g., DVD player for Linux)
    - is it fair? (the entertainment industry wants everyone to pay for their weak copyright protection)
    - MS Palladium (Microsoft's secure computing base) place Microsoft as the gatekeeper of identification and authentication
      - end systems?
      - keys for hardware
  - Secret vs. Public
    - public: integrity protected

Copyright © William C. Cheng

## Practical Use of Keys

- ↳ Email (PEM or S/MIME)
  - hashes and message keys to be distributed and signed
- ↳ Conferencing
  - group key management (discussed later)
- ↳ Authentication (discussed later)
- ↳ SSL (details later)
  - and other "real time" protocols
  - key establishment

Copyright © William C. Cheng

## Recovery from Exposed Keys

- ↳ Revocation lists (CRLs)
  - long lists
  - hard to propagate
- ↳ Lifetime / expiration
  - short life allows assurance of validity at time of issue
- ↳ Realtime Validation
  - Online Certificate Status Protocol (OCSP)
    - privacy concerns? (server knows who you have been communicating with)
  - What about existing messages?

Copyright © William C. Cheng

## Other Key Management Issues

- ↳ Key size vs. data size
  - affects security and usability
- ↳ Reuse of keys
  - multiple users, multiple messages
- ↳ Initial exchange
  - the bootstrap/registration problem
    - Ex: Web
      - ◊ use social security numbers?
      - ◊ use "personal" information?
      - 2002, Princeton admission official improperly logged into Yale website using "personal" info
    - confidentiality vs. authentication
    - sometimes you do not really need authentication

Copyright © William C. Cheng

## Other Key Management Issues (Cont...)

- sometimes you do not really need authentication (cont...)
  - client is often unauthenticated (server often does not know who the client is)
  - long term relationship more important
  - if the "real owner" hasn't complained and this client is paying the bills, this client is probably the "real owner"
- ↳ Security architectures
  - put some security requirements together

Copyright © William C. Cheng

## Security Architectures

- ↳ DSSA (Distributed Systems Security Architecture)
  - around 1987, originally from DEC
  - Ex: how to protect against booting from a CD and access all files on harddrive
  - hardware can checksum OS before loading the OS
    - if no match, don't load it
    - if match, create a certificate, pass it to the OS
  - delegation is the important issue
  - workstation can act as user
    - if given key
    - software can act as developer
    - if checksum validated
    - complete chain needed to assume authority
    - roles provide limits on authority - new sub-principal

Copyright © William C. Cheng

## Security Architectures (Cont...)

- ↳ Microsoft Authenticode
  - ↳ downloadable executables such as Java applets, Windows updates, ActiveX controls uses signed certificates
  - ↳ delegate trust to browser
- ↳ Proxies (also based on delegation)
  - ↳ limits on authority explicitly embedded in proxies
  - ↳ works well with ACL (access control list)
  - ↳ more on proxies in "authorization"