

CS551

Multi-homing in BGP

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<http://merlot.usc.edu/cs551-f12>

Multi-homing

- With multi-homing, a single network has more than one connections to the Internet
- Improves reliability and performance: can accommodate link failure
- bandwidth is sum of links to Internet
- Multiple connections provide *load sharing* but *not* load balancing
- BGP cannot do load balancing

Static Routing May Not Work

- Static routing may send traffic to ISPs 2-n from customer over one link and traffic to ISP1 over the other link
- Lacks flexibility (especially when ISP1 grows and shrink)

Transit vs. Nontransit Services

- ISPn provides transit service to Customern
- ISPn provides non-transit service from another ISP for traffic for its customer

- ### Issues With Multi-homing
- Symmetric routing while conventional wisdom prefers symmetric paths, many are asymmetric
 - Packet re-ordering may trigger TCP's fast retransmit algorithm
 - Other concerns: addressing, DNS, aggregation
 - Note: using BGP in multi-homing situation is not an off-the-shelf use of the protocol

Multi-homing to a Single Provider: Case 1

- Easy solution: no BGP, but use IMUX or Multilink PPP
- Harder solution: use BGP
- makes assumptions about traffic (same amount of prefixes can be reached from both links)

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7

Multi-homing to a Single Provider: Case 2

- For ISP->Customer traffic: use MED in Customer or LOCAL-PREF in ISP
- For Customer->ISP traffic: break-down prefix and advertise different prefixes over different links with default routes
- Good if traffic load to/from prefixes is equal
- if single prefix in Customer, only 1 link will be used for ISP->Customer traffic

Good if traffic load to/from prefixes is equal

- Customer learns full BGP routes and load-shares
- Customer learns full BGP (reordering?)
- R3 alternates links
- For Customer->ISP traffic: LOCAL-PREF in ISP
- use MED in Customer or same as before
- For ISP->Customer traffic, same as before

Good if traffic load to/from prefixes is equal

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8

Multi-homing to a Single Provider: Case 3

- Major issues:
 - addressing
 - aggregation
- Customer address space: (what are the advantages and disadvantages of each approach?)
 - delegated by ISP1
 - delegated by ISP2
 - delegated by ISP1 and ISP2
 - obtained independently

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Multi-homing to a Single Provider: Case 4

- Most reliable approach
- no equipment sharing
- ISP -> Customer: same as case 3
- Customer -> ISP: same as case 2

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10

Case 1: Customer Uses Address Space From One ISP (1 or 2)

- Customer uses address space from ISP1
- ISP1 advertises /16 aggregate
- Customer advertises /24 route to ISP2
- ISP2 relays route to ISP1 and ISP3
- ISP2-3 use /24 route
- ISP1 routes directly
- Problems with traffic load? ("traffic magnet")
- Note: this can actually work well if the relative sizes of the providers have a good match

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11

Multi-homing to a Single Provider: Case 2

- ISP1 aggregates to a /19 at border router to reduce internal tables
- ISP1 still announces /16
- ISP1 hears /24 from ISP2
- ISP1 routes packets for customer to ISP2
- Workaround: ISP1 must inject /24 into I-BGP

Pitfalls

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Multi-homing to a Single Provider: Case 3

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Case 3: Customer Uses Its Own Address Space

Offers the most control, but at the cost of aggregation

Still need to control paths:
 = suppose ISP1 large, ISP2-3 small
 = customer advertises long path to ISP1, but LOCAL-PREF attribute used to override
 = ISP3 learns shorter path from ISP2

Bottom line: no good and general solution for multi-homing to multiple providers

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16

How Can BGP Express the Following Policies:

- 1 will not act as transit to 3 (do not tell anyone about 3)
- 1 will use the blue path for packets destined to 4 and the red for packets destined to 5 (cannot control path, just first hop)
- 2 will not accept packets sourced in 1 (no way)

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18

Route Flap Dampening (Cont...)

[CISCO - Intro to BGP]

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13

Case 2: Customer Uses Address Space From Both ISPs

ISP1 and ISP2 continue to announce aggregates
 = Load sharing depends on traffic to two prefixes
 = Lack of reliability: if ISP1 link goes down, part of customer becomes inaccessible
 = Customer may announce prefixes to both ISPs, but still problems with longest match as in case 1

138.39.1/24
204.70.1/24

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15

How Can BGP Express the Following Policies:

- 1 will not act as transit to 3
- 1 will use the blue path for packets destined to 4 and the red for packets destined to 5
- 2 will not accept packets sourced in 1

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17

Route Flap Dampening

Problem: route flap when a flaky link constantly goes up and down:
 = BGP sessions disappear and reappear
 = routes are withdrawn and re-advertised
 = global effects (does the flap of a butterfly's wing in Brazil set off a tornado in Texas?)
 = BGP was extended to dampen route flaps
 = Associate a penalty with each route
 = increase when route flaps
 = exponentially decay penalty with time
 = When penalty reaches threshold, suppress route
 = must never forget routes

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22

Some BGP Stats (as of 30-Jan-2003)

- ↳ BGP routing table entries: 120,000 [data from Japan (bgp-stats@lists.apnic.net)]
- ↳ prefixes after max aggregation: 76,596
- ↳ Addresses announced: 1,180,368,745
- ↳ 31.8% of available address space announced
- ↳ 57.9% of the allocated address space announced
- ↳ 55.0% of available address space allocated
- ↳ AS's in Internet routing table: 14,513
- ↳ origin-only AS's: 12,615
- ↳ transit AS's: 1,898
- ↳ AS path length
 - ↳ mean: 5.3
 - ↳ maximum seen: 17

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20

BGP Limitations: Policy

↳ Ex: fish routing

```

graph TD
  D --- A
  D --- B
  D --- C
  D --- E
  A --- B
  B --- C
  E --- F
  
```

- ↳ E would like A to send traffic to E via B while
- ↳ F would like A to send traffic to F via C
- ↳ but how?

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21

Other BGP-related Issues

- ↳ Convergence Time [Labovitz0a]
- ↳ Router synchronization [Floyd94b]
- ↳ Congestion [Shaikh00a]
- ↳ Policy and convergence [Gao00a, Tangmunarunkit01a]
- ↳ Misconfiguration [Mahajan02a]
- ↳ Other other issues
 - ↳ routing arbiter - central DB of policies
 - ↳ robustness in the face of router resource exhaustion [Chang, Govindan, Heidemann]

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19

Tricky Issues

- ↳ "Synchronizing" intra and inter-domain routing
- ↳ Getting packets to the right exit router without introducing too much flux into intra-domain routing
- ↳ Multi-homing
 - ↳ interaction with aggregation
 - ↳ How much policy *should* we actually be able to support???