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


Mobile Routing Alternatives

- Why not just assume a base station?
 - good for many cases, but not some (military, disaster recovery, sensor nets)
- Why not just use regular Internet routing?
 - completely different assumptions about stability, hierarchy
 - but might work well for fixed wireless
- Why not just flood the data to all nodes?
 - best choice for rapid movement, but doesn't work for many nodes and much traffic

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Problems With Traditional Approaches

- Periodic (a priori) routing or LS updates are expensive
- Dynamic topology problems:
 - frequent LS updates
 - algorithm must converge very quickly to avoid blackholes
- Many "links" in wireless (many nodes can hear each other)
 - large routing tables/messages
- Not studied in the context of realistic radio propagation models, MAC layers and mobility patterns

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


Problems Using DV or LS

- DV protocols may form loops
 - very wasteful in wireless: bandwidth, power
 - loop avoidance sometimes complex
- LS protocols: high storage and communication overhead
 - More links in wireless (e.g., clusters) - may be redundant
 - higher protocol overhead

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
Ad-hoc Routing

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<http://merlot.usc.edu/cs51-f12>

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Ad-hoc Routing Goals

- Paths should be:
 - multi-hop paths
 - loop-free
 - inexpensive to set up
 - robust to node movement
- System should autoconfigure, adapt to movement
 - Low memory, bandwidth, energy required
 - scalable with numbers of nodes
 - localize effects of link failure
- Multicast?
 - not here :-)

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Ad Hoc Routing

- Create multi-hop connectivity among set of wireless, possibly moving, nodes
 - typically trustworthy (more or less)
 - want low energy consumption
 - implications: cannot use hierarchical addressing, assume things will change
- Mobile, wireless hosts act as forwarding nodes as well as end systems
 - Need routing protocol to find multi-hop paths
 - Needs to be dynamic to adapt to new routes, movement
 - Interesting challenges related to interference and power limitations

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Proposed Protocols

- A priori protocols (DSDV, TORA) pre-compute routing tables
 - = Destination-Sequenced Distance Vector (DSDV)
 - hbh, DV protocol w/periodic routing update broadcasts
 - = Temporally-Ordered Routing Algorithm (TORA)
 - on demand creation of hbh routes based on link-reversal
- On-demand protocols (DSR, AODV) compute routes lazily
 - = *Dynamic Source Routing (DSR)*
 - on demand source route discovery
 - = Ad Hoc On-Demand Distance Vector (AODV)
 - combination of DSR and DSDV: on demand route discovery with hbh routing

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Key Ideas

- Ad-hoc routing: routing in multi-hop networks of mobile nodes
 - = no base stations
- Note that this goes way back: DARPA packet radio in the early '70s

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Route Requests (RREQ) and Replies (RREP)

- A request (w/new hop) is forwarded if:
 - = not a duplicate
 - = node is not the destination (if D, then send RREP reply)
 - = node not already listed in recorded source route (loop suppression)
- Replies at D are returned to S via
 - = source route from D's cache
 - = invert source route (not preferred because of potential asymmetry)
 - = piggy-backed on another route request from D to S

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..Problems

- Periodic updates waste power
 - = tx sends portion of battery power into air
 - = reception requires less power, but periodic updates prevent mobile from "sleeping"
- Convergence may be slower in conventional networks but frequent updates
 - = fast in ad-hoc networks and be done without

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Dynamic Source Routing in Wireless Networks

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
DSR

Components:

- route discovery
- = route maintenance
- Route discovery - basic idea
 - = S floods route-request to D
 - = each node forwards request by adding own address and re-broadcasting
 - = requests propagate outward until target is found

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Route Cache

- Keep all source routes in route cache (reuse if possible)
- Can short-circuit RREQs if it matches the cache
- Just generate a reply
- but make sure you don't step on your neighbors
- Can eavesdrop to pre-fill route cache

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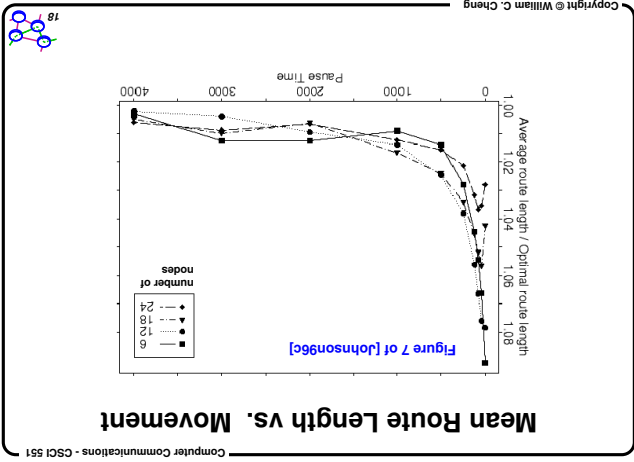
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Sending Data

- Check cache for route to D
- If route exists then
 - If reachable in one hop
 - send packet
 - else insert routing header to D and send
- If route does not exist, buffer packet and initiate route discovery

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Route Maintenance

- Several options to detect broken "links":
 - use link-level info (if present)
 - passive ACKs (listen for data to be forwarded)
 - end-to-end

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Other Optimizations

- Piggybacking
 - Data messages on the initial route request
 - Reply messages on the reverse route request
 - Need to do this carefully
 - ... interacts with route cache optimization
- Hop short-cuts
 - If a node notices that the packet has skipped a hop, it can send an unsolicited route reply
- Optimized error handling
 - Rate limiting requests: need to worry about partitions
 - Snooping error messages

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Performance Evaluation

- Models for
 - traffic: random pairs sending pseudo-CBR
 - mobility: random waypoint
 - node placement: random
- Metrics
 - path-length relative to optimal
 - message count relative to optimal

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Discussion

- ↳ Context = Ad-hoc networks was in its early stages of development
- ↳ Impact = Showed that these networks needed a different routing architecture
- ↳ Discussion = Good, clean design
- ↳ = Simple and useful performance optimization tricks
- ↳ Other comments = The field is still looking for non-military applications

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Discussion

- ↳ Source routing is good for on-demand routes instead of a priori distribution
- ↳ Route discovery protocol used to obtain routes on demand
- ↳ = Caching used to minimize use of discovery
- ↳ Periodic messages avoided
- ↳ But need to buffer packets

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Other Observations?

- ↳ What do people use ad hoc routing today?
 - = sensor nets
 - = cars on highways (mostly single-hop)
 - = military/disaster recovery applications
 - = mesh networking (at places where there is no infrastructure)
- ↳ small-scale sharing with workgroups could use it, but today most people use floppy/zip disks or CD-RW
- ↳ vs. most laptop users today use basestations
- ↳ What are alternatives?

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