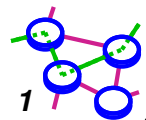


CS551

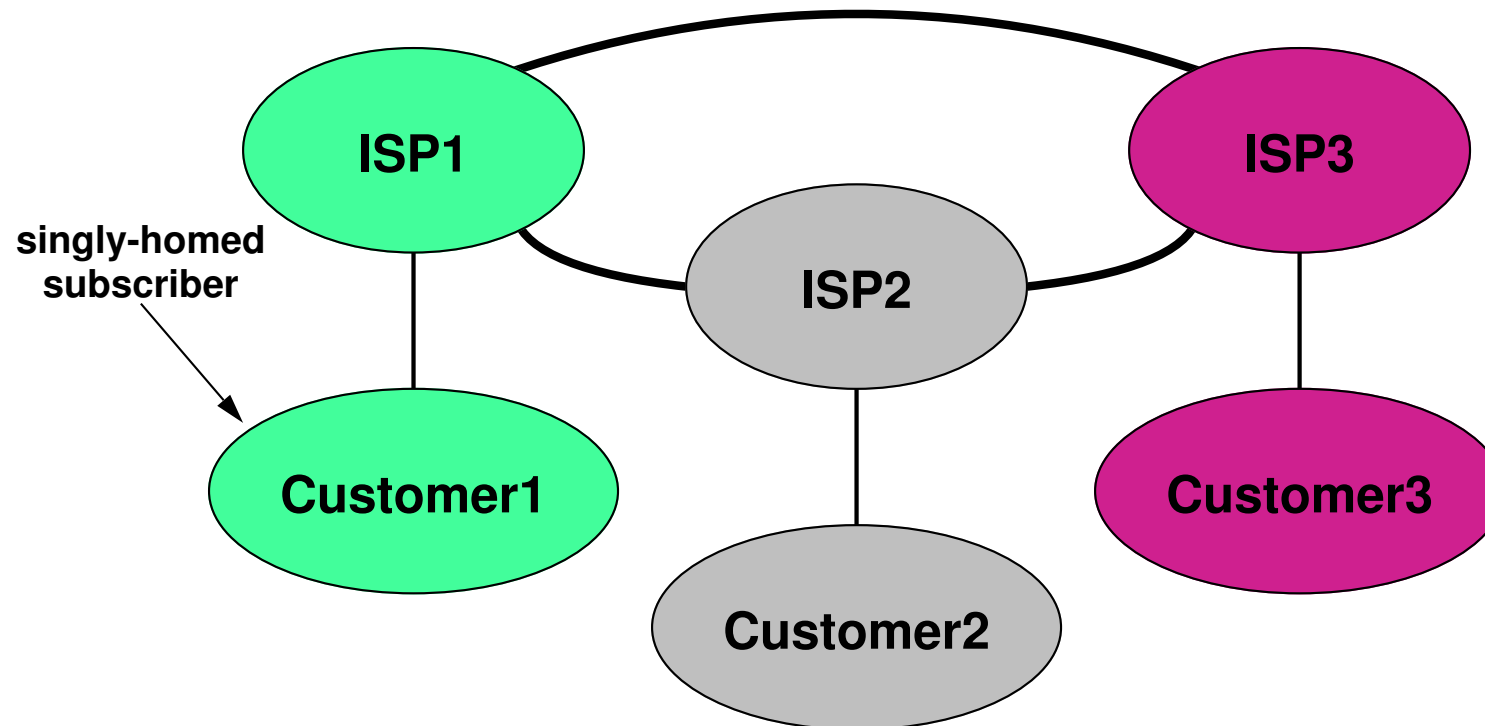
Multi-homing in BGP

Bill Cheng

<http://merlot.usc.edu/cs551-f12>

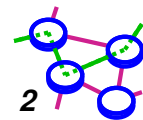


Transit vs. Nontransit Services



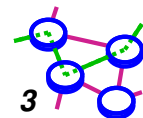
➡ ISP_n provides transit service to $Customer_n$

➡ ISP_n provides non-transit service from another ISP for traffic for its customer



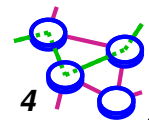
Multi-homing

- ➔ With multi-homing, a single network has more than one connections to the Internet
- ➔ Improves reliability and performance:
 - ▬ can accommodate link failure
 - ▬ bandwidth is sum of links to Internet
- ➔ Multiple connections provide *load sharing* but *not* load balancing
 - ▬ BGP cannot do load balancing

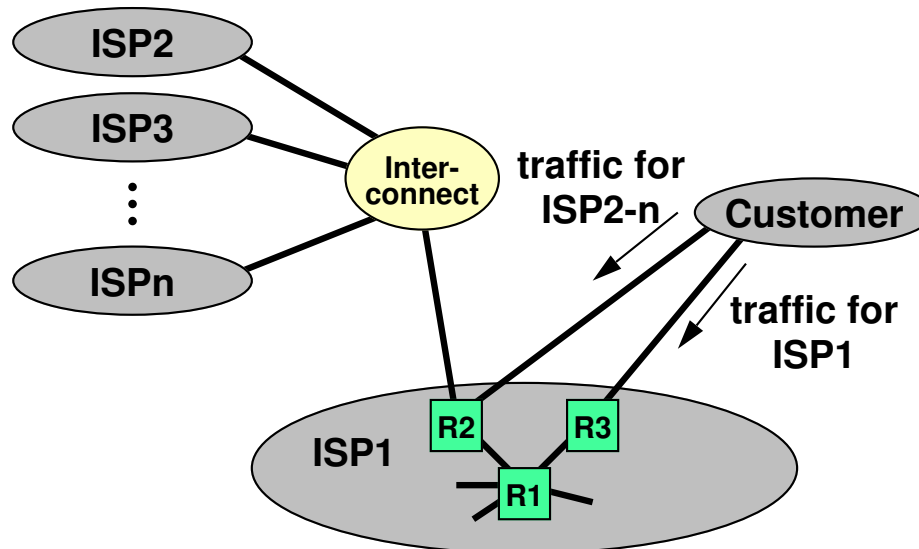


Issues With Multi-homing

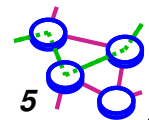
- ➔ **Symmetric routing**
 - ➔ while conventional wisdom prefers symmetric paths, many are asymmetric
- ➔ **Packet re-ordering**
 - ➔ may trigger TCP's fast retransmit algorithm
- ➔ **Other concerns:**
 - ➔ addressing, DNS, aggregation
- ➔ **Note:** using BGP in multi-homing situation is not an off-the-shelf use of the protocol



Static Routing May Not Work

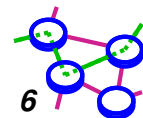
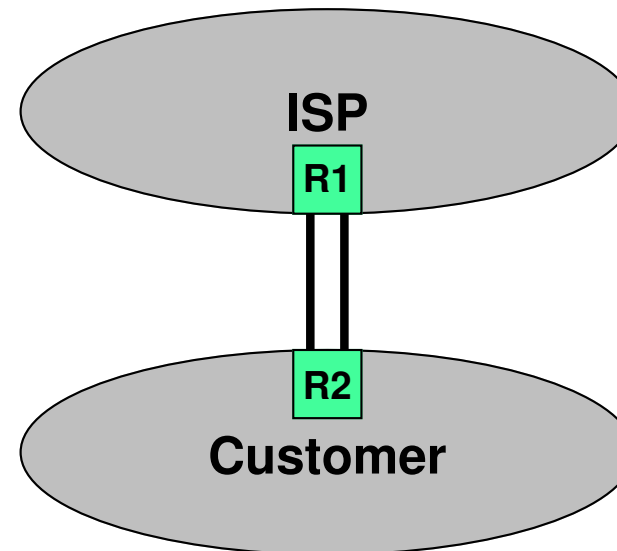


- Static routing may send traffic to ISPs 2-n from customer over one link and traffic to ISP1 over the other link
- Lacks flexibility (especially when ISP1 grows and shrink)



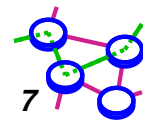
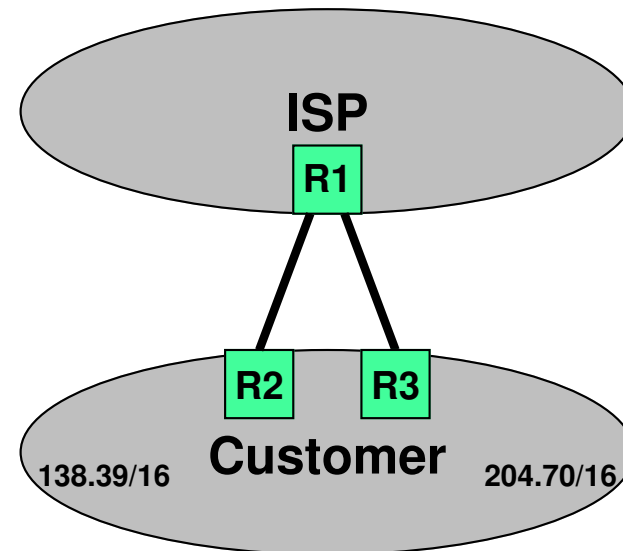
Multi-homing to a Single Provider: Case 1

- ➔ **Easy solution:**
 - ▬ no BGP, but use IMUX or Multilink PPP
- ➔ **Harder solution:**
 - ▬ use BGP
 - ▬ makes assumptions about traffic (same amount of prefixes can be reached from both links)



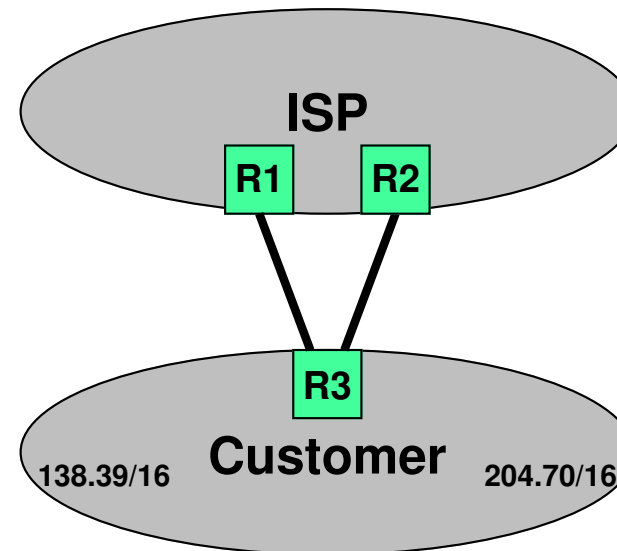
Multi-homing to a Single Provider: Case 2

- ➔ For ISP-> Customer traffic:
 - ▢ use MED in Customer or LOCAL-PREF in ISP
- ➔ For Customer->ISP traffic:
 - ▢ break-down prefix and advertise different prefixes over different links with default routes
- ➔ Good if traffic load to/from prefixes is equal
 - ▢ if single prefix in Customer, only 1 link will be used for ISP->Customer traffic



Multi-homing to a Single Provider: Case 3

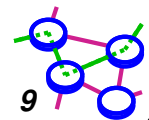
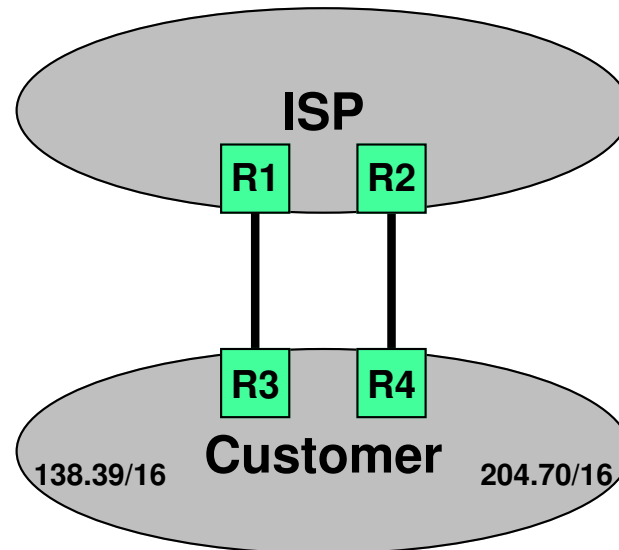
- ➔ For ISP->Customer traffic, same as before:
 - ➔ use MED in Customer or LOCAL-PREF in ISP
- ➔ For Customer->ISP traffic:
 - ➔ R3 alternates links (reordering?)
 - ➔ Customer learns full BGP routes and load-shares



- ➔ Good if traffic load to/from prefixes is equal

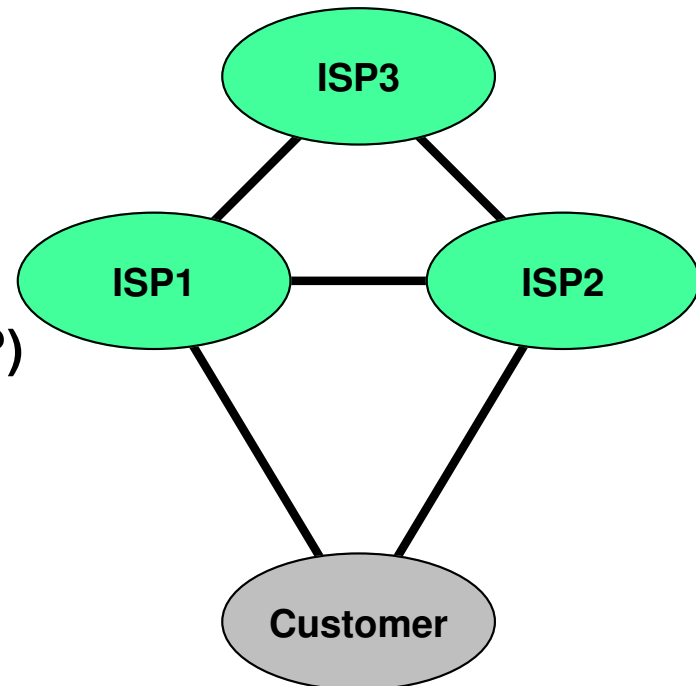
Multi-homing to a Single Provider: Case 4

- ➔ **Most reliable approach**
 - ▢ no equipment sharing
- ➔ **ISP -> Customer:**
 - ▢ same as case 3
- ➔ **Customer -> ISP:**
 - ▢ same as case 2



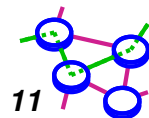
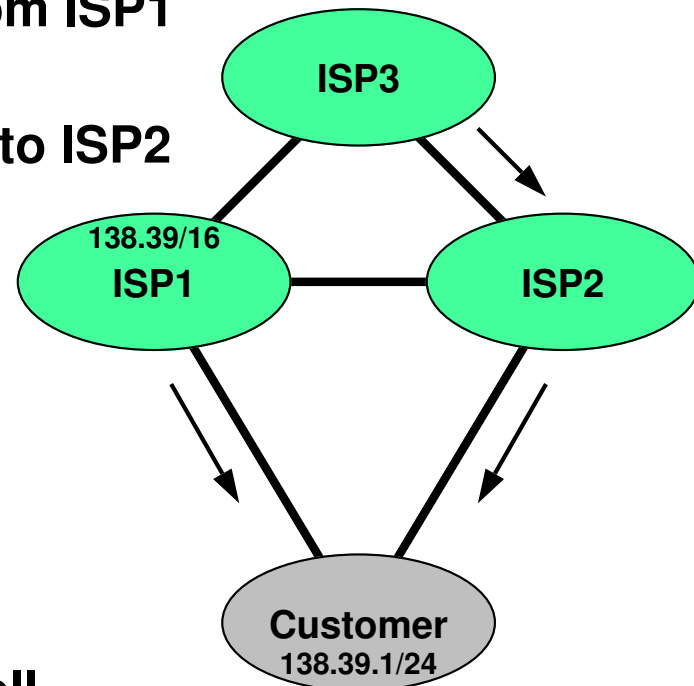
Multi-homing to Multiple Providers

- ➔ Major issues:
- ➔ addressing
 - ➔ aggregation
- ➔ Customer address space:
(what are the advantages and disadvantages of each approach?)
- ➔ delegated by ISP1
 - ➔ delegated by ISP2
 - ➔ delegated by ISP1 and ISP2
 - ➔ obtained independently



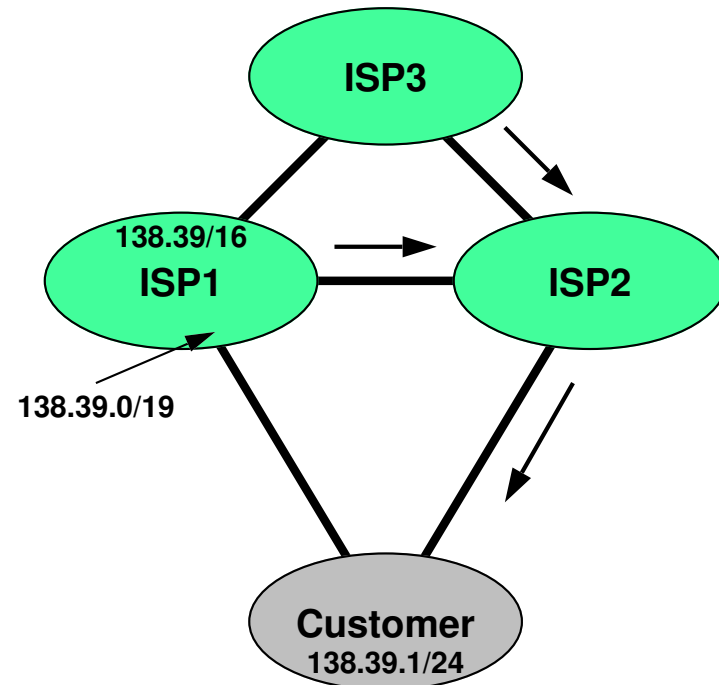
Case 1: Customer Uses Address Space From One ISP (1 or 2)

- ➔ Customer uses address space from ISP1
- ISP1 advertises /16 aggregate
 - Customer advertises /24 route to ISP2
 - ISP2 relays route to ISP1 and ISP3
 - ISP2-3 use /24 route
 - ISP1 routes directly
 - Problems with traffic load? (longest prefix becomes a *"traffic magnet"*)
 - *Note:* this can actually work well if the relative sizes of the providers have a good match



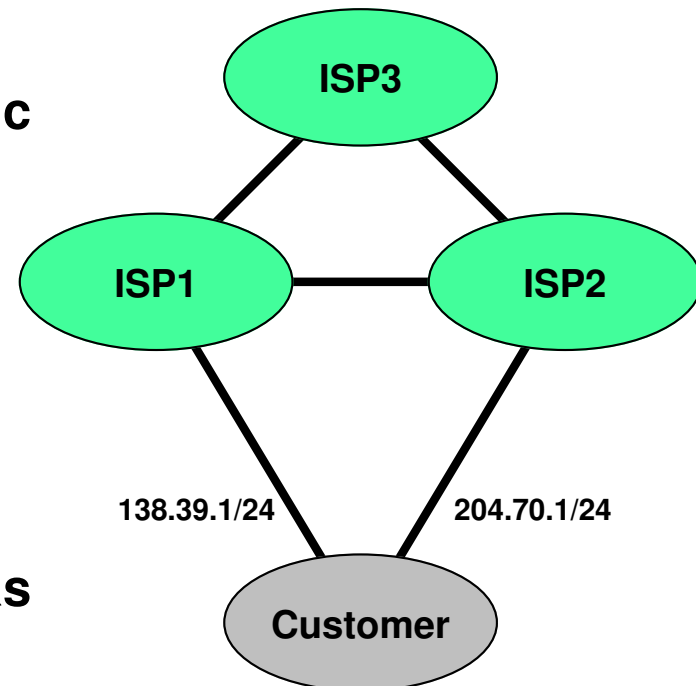
Pitfalls

- ➔ **ISP1 aggregates to a /19 at border router to reduce internal tables**
- ➔ **ISP1 still announces /16**
 - ➔ **ISP1 hears /24 from ISP2**
 - ➔ **ISP1 routes packets for customer to ISP2!**
 - ➔ **Workaround: ISP1 must inject /24 into I-BGP**



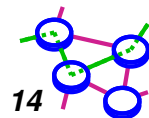
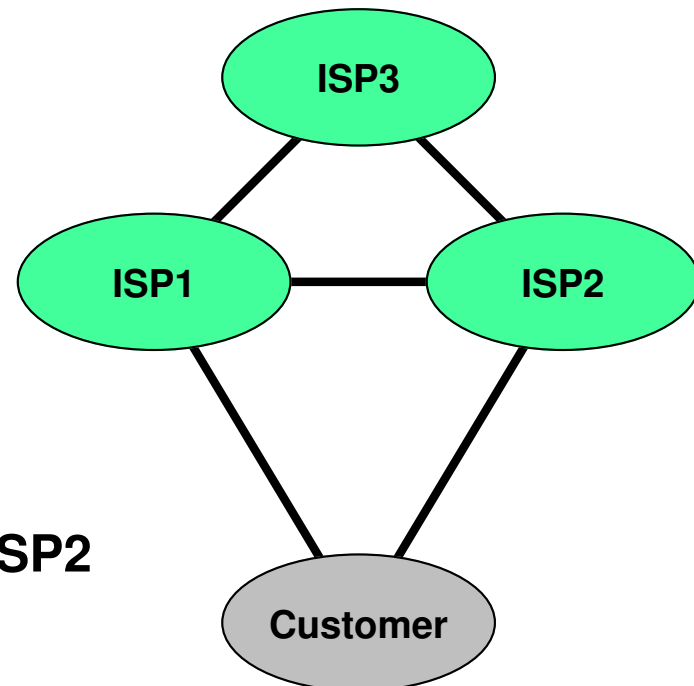
Case 2: Customer Uses Address Space From Both ISPs

- ➡ ISP1 and ISP2 continue to announce aggregates
- ➡ Load sharing depends on traffic to two prefixes
 - ➡ Lack of reliability: if ISP1 link goes down, part of customer becomes inaccessible
 - ➡ Customer may announce prefixes to both ISPs, but still problems with longest match as in case 1



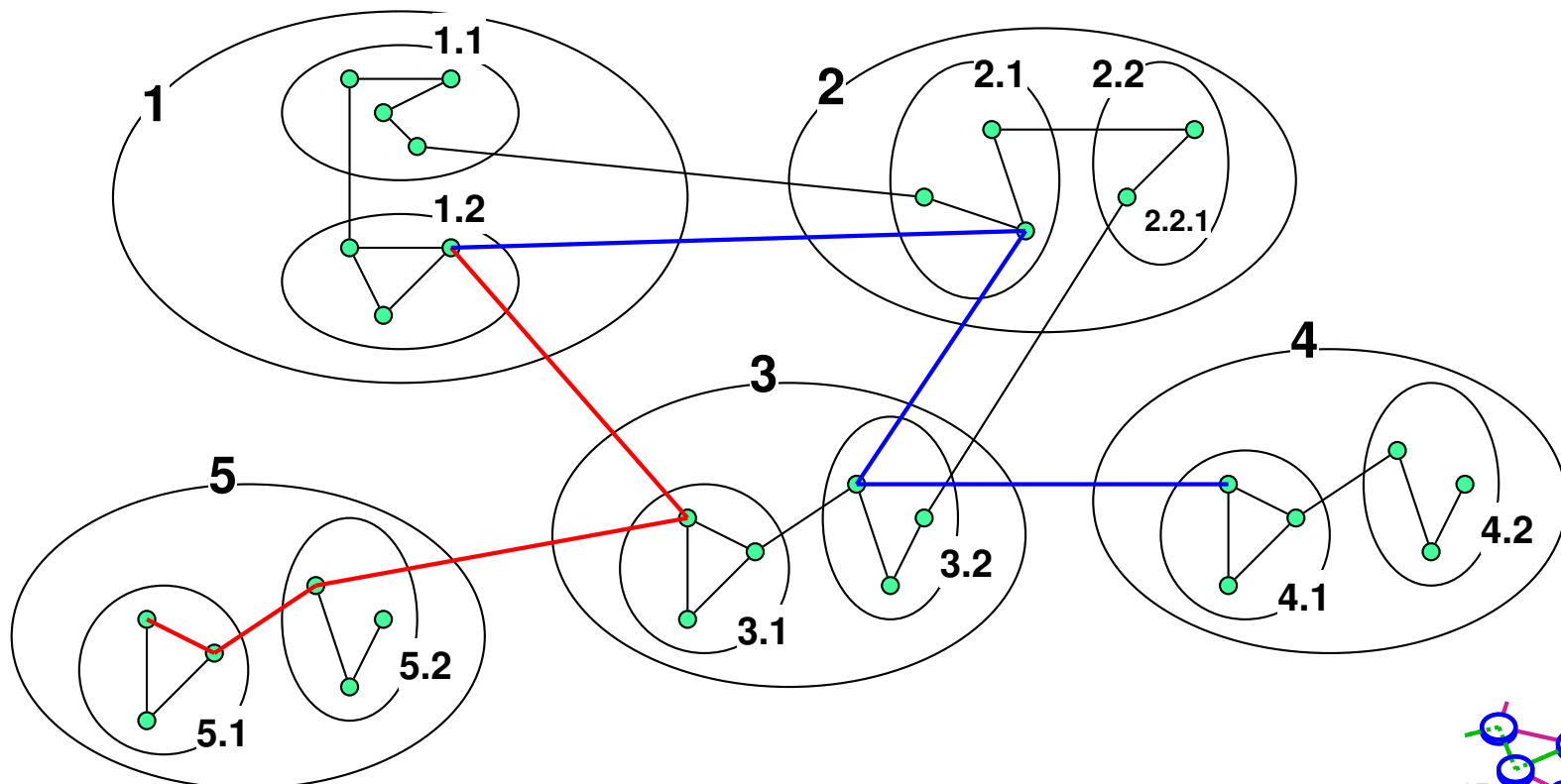
Case 3: Customer Uses Its Own Address Space

- ➔ Offers the most control, but at the cost of aggregation
- ➔ Still need to control paths:
 - suppose ISP1 large, ISP2-3 small
 - customer advertises long path to ISP1, but LOCAL-PREF attribute used to override
 - ISP3 learns shorter path from ISP2
- ➔ **Bottom line:** no good and general solution for multi-homing to multiple providers



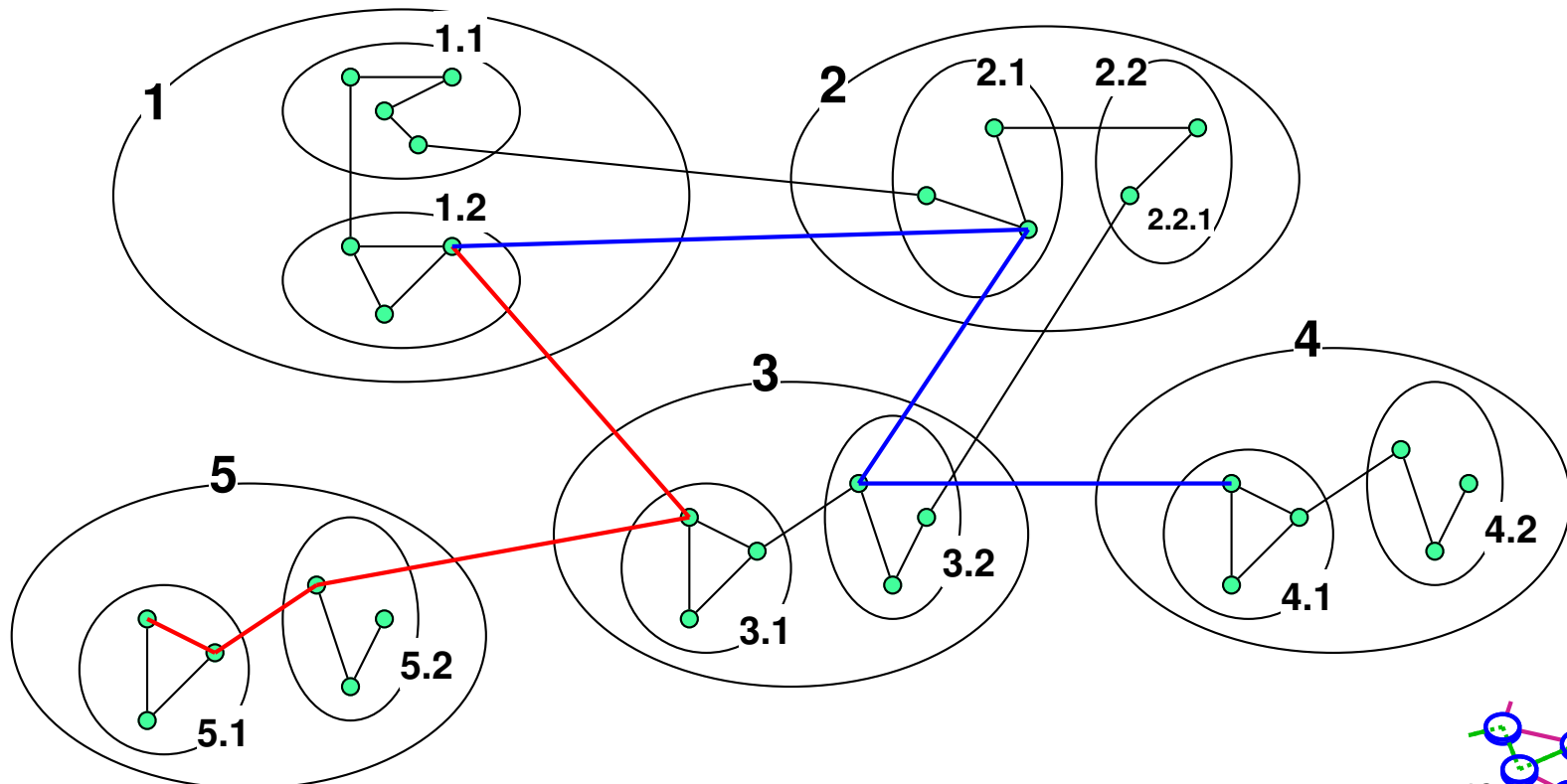
How Can BGP Express the Following Policies:

- ❑ 2 will not act as transit to 3
- ❑ 1 will use the blue path for packets destined to 4 and the red for packets destined to 5
- ❑ 2 will not accept packets sourced in 1



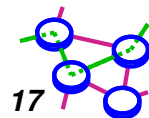
How Can BGP Express the Following Policies:

- ❑ 2 will not act as transit to 3 (*do not tell anyone about 3*)
- ❑ 1 will use the blue path for packets destined to 4 and the red for packets destined to 5 (*cannot control path, just first hop*)
- ❑ 2 will not accept packets sourced in 1 (*no way*)



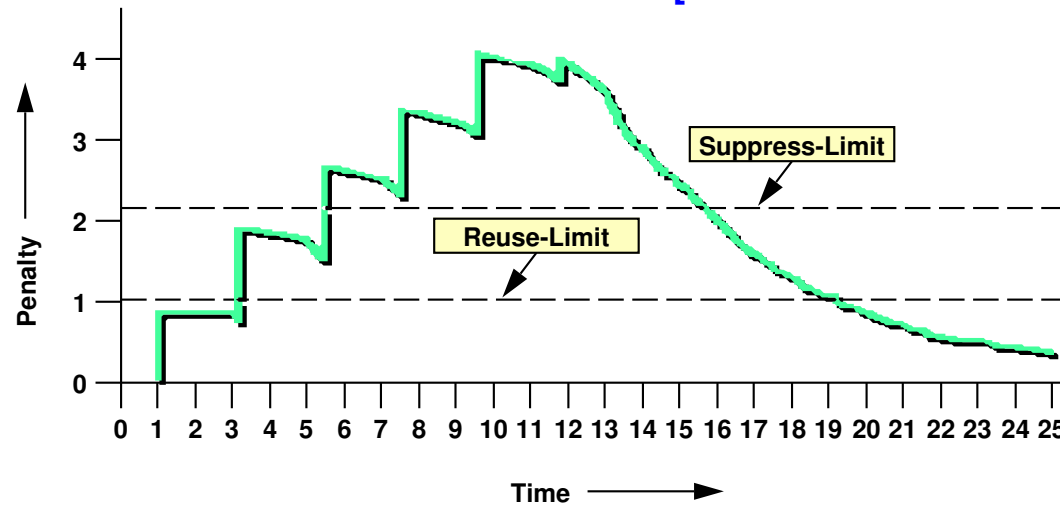
Route Flap Dampening

- ➔ **Problem: route flap when a flaky link constantly goes up and down:**
 - ➔ BGP sessions disappear and reappear
 - ➔ routes are withdrawn and re-advertised
 - ➔ global effects (does the flap of a butterfly's wing in Brazil set off a tornado in Texas?)
- ➔ **BGP was extended to dampen route flaps**
- ➔ **Associate a penalty with each route**
 - ➔ increase when route flaps
 - ➔ exponentially decay penalty with time
- ➔ **When penalty reaches threshold, suppress route**
 - ➔ must never forget routes



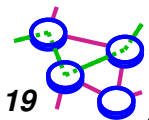
Route Flap Dampening (Cont...)

[CISCO - Intro to BGP]



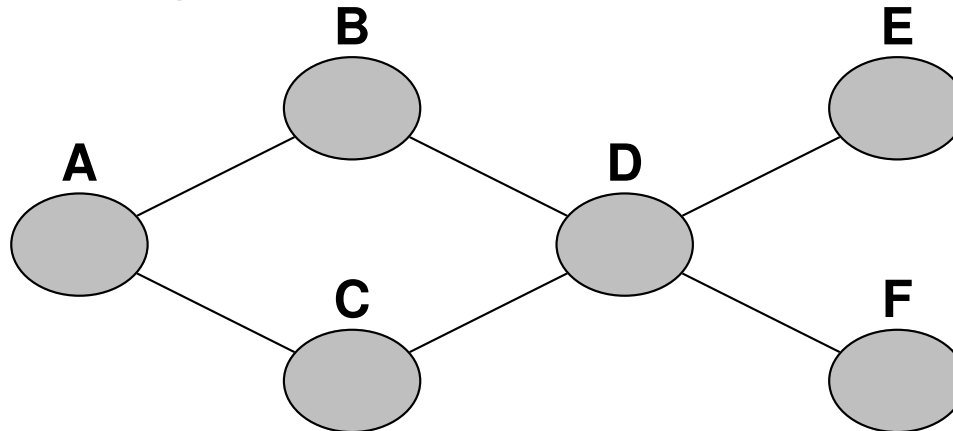
Tricky Issues

- ➔ "Synchronizing" intra and inter-domain routing
- ➔ Getting packets to the right exit router without introducing too much flux into intra-domain routing
- ➔ Multi-homing
 - ▬ interaction with aggregation
- ➔ How much policy *should* we actually be able to support???



BGP Limitations: Policy

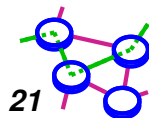
➔ Ex: *fish routing*



- ➔ E would like A to send traffic to E via B while F would like A to send traffic to F via C
- ➔ but how?

Other BGP-related Issues

- ➔ **Convergence Time [Labovitz00a]**
- ➔ **Router synchronization [Floyd94b]**
- ➔ **Congestion [Shaikh00a]**
- ➔ **Policy and convergence [Gao00a, Tangmunarunkit01a]**
- ➔ **Misconfiguration [Mahajan02a]**
- ➔ **Other other issues**
 - ▬ **routing arbiter - central DB of policies**
 - ▬ **robustness in the face of router resource exhaustion [Chang, Govindan, Heidemann]**



Some BGP Stats (as of 30-Jan-2003)

[data from Japan

(bgp-stats@lists.apnic.net)]

- ➔ BGP routing table entries: 120,000
 - ▬ prefixes after max aggregation: 76,596

- ➔ Addresses announced: 1,180,368,745
 - ▬ 31.8% of available address space announced
 - ▬ 57.9% of the allocated address space announced
 - ▬ 55.0% of available address space allocated

- ➔ AS's in Internet routing table: 14,513
 - ▬ origin-only AS's: 12,615
 - origin-only AS's with only 1 prefix: 5,690
 - ▬ transit AS's: 1,898

- ➔ AS path length
 - ▬ mean: 5.3
 - ▬ maximum seen: 17

