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### Multi-threading Exercise

- Make sure you are familiar with the [pthreads](#) library
- good source is the book by Nichols, Buttar, and Farrell *"Pthreads Programming"*, O'Reilly & Associates, 1996
- you must learn how to use mutex and condition variables correctly
- `pthread_mutex_lock()`/`pthread_mutex_unlock()`
- `pthread_cond_wait()`/`pthread_cond_signal()`
- `pthread_cond_broadcast()`
- you must learn how to handle UNIX signals
- `pthread_sigmask()`/`sigwait()`
- `pthread_cancelstate()`
- `pthread_setcanceltype()`
- `pthread_testcancel()`

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### pthread\_sigmask()

Child thread example

- child thread unblocks SIGINT
- child thread is designated to handle SIGINT, no other

```

struct sigaction act;
void * handler(char argv[]);
act.sa_handler = interrupt;
sigaction(SIGINT, &act, NULL);
pthread_sigmask(SIG_UNBLOCK, &new, NULL);
printf("\n Press CTRL-C to deliver SIGINT\n");
sleep(8); /* give user time to hit CTRL-C */
void interrupt(int sig)
{
    printf("thread %d caught signal %d\n", thr_self(), sig);
}
    
```

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### Arrivals & Departures

- $a_i$ : arrival time
- $d_i$ : departure time
- $s_i$ : service time
- $r_i$ : response (system) time
- $q_i$ : queuing time

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### Queueing Abstraction

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### pthread\_sigmask()

Look at the man pages of `pthread_sigmask()` on nunki and try to understand the example there

- designate child thread to handle SIGINT
- parent thread blocks SIGINT

```

#include <pthread.h>
pthread_t new;
void * handler(), interrupt();
main(int argc, char *argv[] )
{
    sigmaskset(&new, SIGINT);
    pthread_sigmask(SIG_BLOCK, &new, NULL);
    pthread_join(user_thread, NULL);
    pthread_create(&user_thread, NULL, handler, argv[1]);
    printf("thread handler, %d exited\n", user_thread);
    sleep(2);
}
    
```

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### Ex:

- line at a bank
- multi-processor executing jobs from a shared job queue
- ER

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### Warm-up Project #2

# CS51

## Bill Cheng

<http://merlot.usc.edu/csc51-f12>

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### Arrivals & Departures (Cont...)

$\lambda$   
 $q_1, q_2, q_3 \sim 0$   
 $q_4 > 0$

An **event queue** is a sorted list of events according to timestamps; smallest timestamp at the head of queue  
**Object oriented:** every object has a "next event" (what it will do next if there is no interference), this event is inserted into the event queue  
 Execution: remove an event from the head of queue, "execute" the event (notify the corresponding object so it can insert the next event)  
 Insert into the event queue according to timestamp of a new event; insertion may cause additional events to be deleted or inserted  
 Potentially repeatable runs (if the same seed is used to initialize random number generator)

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### Event Driven Simulation (Cont...)

$\min(a_3, d_1, d_2) = d_1$ , next event to fire is  $[d_1, \text{destroy}(C_1)]$   
 $A : [a_3, \text{create}(C_3)]$     $S1 : \text{NULL}$     $S2 : \text{NULL}$   
 $Q1 : \text{empty}$   
 $\min(a_3, d_2) = a_3$ , next event to fire is  $[a_3, \text{create}(C_3)]$   
 $\text{create}(C_3), Q1 \rightarrow \text{enqueue}(C_3)$   
 $\text{create}(C_3), Q1 \rightarrow \text{serve}(C_3), S1 \rightarrow \text{serve}(C_3)$   
 $A : [a_4, \text{create}(C_4)]$     $S1 : [d_3, \text{destroy}(C_3)]$   
 $Q1 : \text{empty}$     $S2 : [d_2, \text{destroy}(C_2)]$   
 $\min(a_4, d_2, d_3) = a_4$ , next event to fire is  $[a_4, \text{create}(C_4)]$   
 $\text{create}(C_4), Q1 \rightarrow \text{enqueue}(C_4)$   
 $A : [a_5, \text{create}(C_5)]$     $S1 : [d_3, \text{destroy}(C_3)]$     $S2 : [d_2, \text{destroy}(C_2)]$   
 $Q1 : \text{empty}$     $S2 : [d_2, \text{destroy}(C_2)]$   
 etc.

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### Time Driven Simulation

Every active object is a thread  
 = a customer is a passive object; it gets passed around  
 To execute a job for  $x$  msec, the thread sleeps for  $x$  msec  
 = `nunki.usc.edu` does not run a realtime OS  
 = it may not get woken up more than  $x$  msec later, and sometimes, **a lot more** than  $x$  msec later  
 = you need to decide if the extra delay is reasonable or it is due to a bug in your code  
 Let your machine decide which thread to run next (irreproducible results)  
 Compete for resources (such as  $Q1$ ), must use mutex

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### Event Driven Simulation (Cont...)

Ex: 4 objects,  $A$  (arrival),  $Q1$  (passive object, does not generate events),  $S1, S2$   
 = Initially:  
 $A : [a_1, \text{create}(C_1)]$     $S1 : \text{NULL}$     $S2 : \text{NULL}$   
 $Q1 : \text{empty}$   
 = only one event, next event to fire is  $[a_1, \text{create}(C_1)]$   
 $\text{create}(C_1), Q1 \rightarrow \text{enqueue}(C_1)$   
 $Q1 \rightarrow \text{dequeue}(C_1), S1 \rightarrow \text{serve}(C_1)$   
 $A : [a_2, \text{create}(C_2)]$     $S1 : [d_1, \text{destroy}(C_1)]$   
 $Q1 : \text{empty}$     $S2 : \text{NULL}$   
 $\min(a_2, d_1) = a_2$ , next event to fire is  $[a_2, \text{create}(C_2)]$   
 $\text{create}(C_2), Q1 \rightarrow \text{enqueue}(C_2)$   
 $Q1 \rightarrow \text{dequeue}(C_2), S2 \rightarrow \text{serve}(C_2)$   
 $A : [a_3, \text{create}(C_3)]$     $S1 : [d_1, \text{destroy}(C_1)]$     $S2 : \text{empty}$   
 $Q1 : \text{empty}$     $S2 : [d_2, \text{destroy}(C_2)]$

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### Event Driven Simulation (Cont...)

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### Event Driven Simulation (Cont...)

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You will need to implement 3 threads (or 1 main thread and 3 child threads)

- the arrival thread sits in a loop
- sleeps for an interval, trying to match a given interarrival time (from trace or coin flip)
- wakes up, creates a customer object, enqueues the customer to Q1, and goes back to sleep
- if the Q1 was empty before, need to signal or broadcast a queue-not-empty condition
- two server threads
- initially blocked, waiting for the queue-not-empty condition to be signaled
- (cont...)

Time Driven Simulation (Cont...)

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Time Driven Simulation (Cont...)

Notation:  $\alpha_i$ : inter-arrival time for customer  $i$  ( $a_1, a_2, \dots, a_n=0$ )

$\beta_i$ : service time of customer  $i$

Initially:

- A: sleep( $\alpha_1$ )
- Q1: empty
- S1: idle

A wakes up at  $a_1$ : create( $C_1$ ), Q1->enqueue( $C_1$ )

- S1: sleep( $\beta_1$ )
- S2: idle

A wakes up at  $a_1+\alpha_2$ : create( $C_2$ ), Q1->enqueue( $C_2$ )

- S1: sleeping...
- S2: sleep( $\beta_2$ )

etc.

Time Driven Simulation (Cont...)

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Time Driven Simulation (Cont...)

Q: What were we doing when we "added them up"?

A: We were doing "integration"

Hint:  $0 \leq F(x) \leq 1$  for any  $x$

can numerically compute  $F(x)$

$r = \text{drand48}()$

$M = I$

Time Driven Simulation (Cont...)

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Time Driven Simulation (Cont...)

How do you flip a coin according to this distribution?

Think about discrete case:

Add them up:

$r = \text{drand48}() * 23$

lies between 3 and 13, so we have randomly chosen bucket #2

Time Driven Simulation (Cont...)

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Time Driven Simulation (Cont...)

Uniform distribution (pmf), denoted by  $f(x)$

$f(x) = \begin{cases} 1 & 0 \leq x \leq 1 \\ 0 & \text{otherwise} \end{cases}$

$\int_{-\infty}^{\infty} f(x) dx = 1$

Probability Distribution Function (PDF), denoted by  $F(x)$

$F(x) = \int_{-\infty}^x f(w) dw$

$F(x) = \begin{cases} 0 & x < 0 \\ x & 0 \leq x \leq 1 \\ 1 & x > 1 \end{cases}$

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Time Driven Simulation (Cont...)

Q: What were we doing when we "added them up"?

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Hint:  $0 \leq F(x) \leq 1$  for any  $x$

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Time Driven Simulation (Cont...)

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### Calculating Statistics

lock & unlock stdout to print arrival msg  
 arrival thread timeout (read clock)  
 try lock mutex to enter Q  
 enter Q (read clock)  
 unlock mutex  
 lock & unlock stdout to print enter queue msg  
 time in Q

try lock mutex to leave Q  
 leave Q (read clock)  
 unlock mutex  
 lock & unlock stdout to print leave queue msg  
 time in server

begin service  
 leave server  
 charge to no one

select () ?

lock & unlock stdout to print msg  
 time between begin service and leave server is the amount of time in select ()

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### Mean and Standard Deviation

- Average time  
 = for  $n$  samples, add up all the time and divide by  $n$
- Average number of customer at a server  
 = same a fraction of time the server is busy
- Average number of customer at Q1
- Standard deviation is the square root of variance  

$$Var[X] = E[X^2] - (E[X])^2$$

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### Coin Flipping (Cont..)

Note: Inter-arrival time of a Poisson process is Exponentially distributed

Exponential distribution

$f(x) = me^{-mx}$

$f(x) = \int_0^x f(y)dy = 1 - e^{-mx}$

$f = \text{drand48}()$